## B: Boggle Sort

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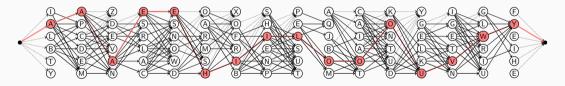
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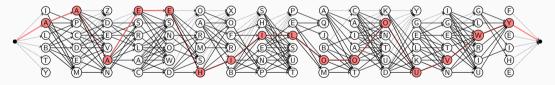
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Brute-force-y sol.: Among the 4 sideways faces, it is optimal to take the alphabetically earliest face that still fits. Thus, there are really only 3 different choices per die: keep, turn on smallest side, turn bottom up. This gives  $3^{16}=43\,046\,721$  choices, which maybe can be systematically checked.

**Graph-y solution:** Create digraph with vertex set  $6\times 16$ ; connect (r,c) to (r',c+1) if the character in (line l, column c) precedes the character in (line r', column c+1) in the alphabet. The weight is 0 if r=1 (dotted), 2 if r=6 (fat), and 1 otherwise. Connect s to (r,1) and (r,16) to t. Then a minimum-weight s, t-path is the solution. Optimal solution for Sample 1:

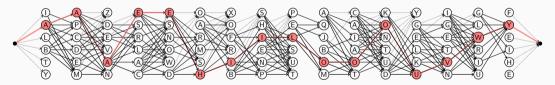


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**Running time:** Using Dijkstra's algorithm, time  $\mathcal{O}(rc \log rc)$  for c dice with r faces. But graph is acyclic, so actually  $\mathcal{O}(rc)$ .

**DP:** Compute, for  $1 \le i \le 16$ , and each letter x, the smallest number f(i,x) of turns needed to bring the first i dice into nondecreasing order such that the ith die shows x. Then, if x appears on the ith die, we have the general case

$$f(i,x) = \max_{y \le x} f(i-1,y)$$

where y ranges over all letters appearing on die (i-1). (Remember the Q=QU pitfall.)

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**Running time:** O(rc) for c dice with r faces.

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**Running time:**  $\mathcal{O}(rc)$  for c dice with r faces.

Statistics: 89 submissions, 35 accepted, 26 unknown